

SAHLQVIST-VAN BENTHEM ALGORITHM

BASED ON THE NOTES BY IAN HODKINSON AND SECTIONS 3.5-3.6 IN BRV

A *boxed atom* is a modal formula of the form $\Box^n p$, for some $n \in \mathbb{N}$, where p is a propositional variable, and $\Box^n p$ is defined by the rule: $\Box^0 p = p$, $\Box^1 p = \Box p$, $\Box^{n+1} p = \Box(\Box^n p)$, $n \in \mathbb{N}$.

A *simple Sahlqvist antecedent* is built from \perp, \top and boxed atoms by applying \diamond and \wedge .

A *simple Sahlqvist formula* is a modal formula of the form $\varphi \rightarrow \psi$, where φ is a simple Sahlqvist antecedent and ψ is a positive formula.

A *Sahlqvist antecedent* is built from \perp, \top , negative formulas and boxed atoms by applying \diamond and \wedge .

A *Sahlqvist implication* is a modal formula of the form $\varphi \rightarrow \psi$, where φ is a Sahlqvist antecedent and ψ is a positive formula.

A *Sahlqvist formula* is built from Sahlqvist implications by applying \Box and \vee .

Theorem (Sahlqvist correspondence) For any Sahlqvist formula φ , there is a corresponding first-order sentence that holds in a frame iff φ is valid in the frame.

This sentence can be obtained from φ by a simple Sahlqvist-van Benthem algorithm. For simplicity we will consider only the case of simple Sahlqvist formulas.

Let φ be a simple Sahlqvist formula.

- (1) Identify boxed atoms in the antecedent.
- (2) Draw the picture that discusses the minimal valuation that makes the antecedent true. Name the worlds involved by t_0, \dots, t_n .
- (3) Work out the minimal valuation i.e., get a first-order expression for it in terms of the named worlds.
- (4) Work out the standard translation of φ . Use the names you fixed for the variables that correspond to \diamond 's in the antecedent.
- (5) Pull out the quantifiers that bind t_i variables in the antecedent to the front. For this use the equivalences

$$\exists x \alpha(x) \wedge \beta \leftrightarrow \exists x (\alpha(x) \wedge \beta),$$

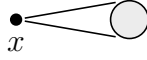
$$\exists x \alpha(x) \rightarrow \beta \leftrightarrow \forall x (\alpha(x) \rightarrow \beta).$$

- (6) Replace all the predicates $P(x), Q(x)$, etc., with the first-order expression corresponding to the minimal valuation.
- (7) Simplify, if possible.
- (8) Add $\forall x$ (binding the free variable of the standard translation) to the resulting first-order formula to obtain the global first-order correspondent.

We will look at a few examples.

Let $\varphi = \Box p \rightarrow p$.

The diagram:



The minimal valuation is $V(p) = \{z : Rxz\}$.

The standard translation of φ is $\forall y (Rxy \rightarrow P(y)) \rightarrow P(x)$.

Replace $P(y)$ with Rxy and $P(x)$ with Rxx . We obtain $\forall y (Rxy \rightarrow Rxy) \rightarrow Rxx$.

This is equivalent to Rxx . By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x Rxx \quad \text{reflexivity!}$$

Let $\varphi = \Box p \rightarrow \Box \Box p$.

The diagram:



The minimal valuation is $V(p) = \{z : Rxz\}$.

The standard translation of φ is

$$\forall y (Rxy \rightarrow P(y)) \rightarrow \forall z (Rxz \rightarrow \forall u (Rzu \rightarrow P(u)))$$

Replace $P(y)$ with Rxy and $P(u)$ with Rxu . We obtain

$$\forall y (Rxy \rightarrow Rxy) \rightarrow \forall z (Rxz \rightarrow \forall u (Rzu \rightarrow Rxu))$$

This is equivalent to

$$\forall z (Rxz \rightarrow \forall u (Rzu \rightarrow Rxu))$$

which is equivalent to

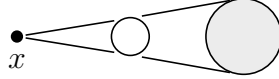
$$\forall z \forall u (Rxz \wedge Rzu \rightarrow Rxu)$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall z \forall u (Rxz \wedge Rzu \rightarrow Rxu) \quad \textbf{transitivity!}$$

Let $\varphi = \Box \Box p \rightarrow \Box p$.

The diagram:



The minimal valuation is $V(p) = \{z : \exists v (Rzv \wedge Rvz)\}$. The standard translation of φ is

$$\forall y (Rxy \rightarrow \forall z (Ryz \rightarrow P(z))) \rightarrow \forall u (Rxu \rightarrow P(u))$$

Replace $P(u)$ with $\exists v (Rzv \wedge Rvz)$. In the antecedent we can replace $P(z)$ with the minimal valuation, but let us note that the instantiation of the standard translation of boxed atoms always gives us a tautology.

We obtain

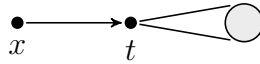
$$\forall u (Rxu \rightarrow \exists v (Rzv \wedge Rvz))$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall u (Rxu \rightarrow \exists v (Rzv \wedge Rvz)) \quad \textbf{density!}$$

Let $\varphi = \Diamond \Box p \rightarrow p$.

The diagram:



The minimal valuation is $V(p) = \{z : Rtz\}$.

The standard translation of φ is

$$\exists t (Rxt \wedge \forall z (Rtz \rightarrow P(z))) \rightarrow P(x)$$

Pull out the existential quantifier in the antecedent. We obtain

$$\forall t (Rxt \wedge \forall z (Rtz \rightarrow P(z))) \rightarrow P(x)$$

Replace $P(z)$ with Rtz and $P(x)$ with Rtx . We obtain

$$\forall t(Rxt \wedge \forall z(Rtz \rightarrow Rtz) \rightarrow Rtx)$$

This is equivalent to

$$\forall t(Rxt \rightarrow Rtx)$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall t(Rxt \rightarrow Rtx) \quad \text{symmetry!}$$

Let $\varphi = p \rightarrow \Diamond p$.

The diagram:



The minimal valuation is $V(p) = \{z : x = z\}$.

The standard translation of φ is

$$P(x) \rightarrow \exists y(Rty \wedge P(y))$$

Replace $P(y)$ with $x = y$ and note that the instantiation of the standard translation of boxed atoms is a tautology. We obtain

$$\exists y(Rxy \wedge y = x)$$

This is equivalent to

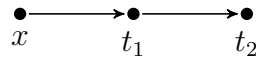
$$Rxx$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x Rxx \quad \text{reflexivity!}$$

Let $\varphi = \Diamond \Diamond p \rightarrow \Diamond p$.

The diagram:



The minimal valuation is $V(p) = \{z : t_2 = z\}$.

The standard translation of φ is

$$\exists t_1(Rxt_1 \wedge \exists t_2(Rt_1t_2 \wedge P(t_2))) \rightarrow \exists y(Rxy \wedge P(y))$$

Pull out the existential quantifiers in the antecedent. We obtain

$$\forall t_1 \forall t_2 ((Rxt_1 \wedge Rt_1t_2 \wedge P(t_2)) \rightarrow \exists y(Rxy \wedge P(y)))$$

Replace $P(y)$ with $t_2 = y$ and note that the instantiation of the standard translation of boxed atoms is a tautology. We obtain

$$\forall t_1 \forall t_2 ((Rxt_1 \wedge Rt_1 t_2) \rightarrow \exists y (Rxy \wedge (y = t_2)))$$

This is equivalent to

$$\forall t_1 \forall t_2 ((Rxt_1 \wedge Rt_1 t_2) \rightarrow Rxt_2)$$

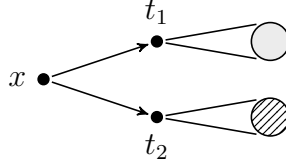
By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall t_1 \forall t_2 ((Rxt_1 \wedge Rt_1 t_2) \rightarrow Rxt_2) \quad \textbf{transitivity!}$$

If φ is a Sahlqvist formula, say $\Box(\varphi \rightarrow \psi) \vee \Box(\varphi' \rightarrow \psi')$ (where $\varphi \rightarrow \psi$ and $\varphi' \rightarrow \psi'$ are simple Sahlqvist formulas), then draw a diagram where outer \Box 's are treated as \Diamond 's of simple Sahlqvist formulas and \vee is treated as \wedge of simple Sahlqvist formulas.

Let $\varphi = \Box(\Box p \rightarrow q) \vee \Box(\Box q \rightarrow p)$.

The diagram:



The minimal valuation is $V(p) = \{z : Rt_1 z\}$ and $V(q) = \{z : Rt_2 z\}$.

The standard translation of φ (keeping in mind t_1 and t_2) is

$$\forall t_1 (Rxt_1 \rightarrow (ST_{t_1}(\Box p) \rightarrow Q(t_1))) \vee \forall t_2 (Rxt_2 \rightarrow (ST_{t_2}(\Box q) \rightarrow P(t_2)))$$

Pull out the quantifiers and replace $Q(t_1)$ with $Rt_2 t_1$ and $P(t_2)$ with $Rt_1 t_2$. Note again that the instantiation of the standard translation of boxed atoms gives a tautology.

We obtain

$$\forall t_1 \forall t_2 ((Rxt_1 \rightarrow Rt_2 t_1) \vee (Rxt_2 \rightarrow Rt_2 t_1))$$

which is equivalent to

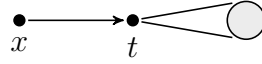
$$\forall t_1 \forall t_2 ((Rxt_1 \wedge Rxt_2) \rightarrow (Rt_1 t_2 \vee Rt_2 t_1))$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall t_1 \forall t_2 ((Rxt_1 \wedge Rxt_2) \rightarrow (Rt_1 t_2 \vee Rt_2 t_1)) \quad \textbf{linearity!}$$

Let $\varphi = \Box(\Box p \rightarrow p)$.

The diagram:



The minimal valuation is $V(p) = \{z : Rtz\}$.

The standard translation of φ (keeping in mind t) is

$$\forall t(Rxt \rightarrow (ST_t(\Box p) \rightarrow P(t)))$$

Replace $P(t)$ with Rtt . Note again that the instantiation of the standard translation of boxed atoms gives a tautology.

We obtain

$$\forall t(Rxt \rightarrow Rtt)$$

By adding $\forall x$ we obtain the global first-order correspondent

$$\forall x \forall t(Rxt \rightarrow Rtt) \text{ every successor is reflexive!}$$